

Internship proposal for Tristan Charrier

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Most multiprocessors exhibit subtle weak memory behaviour, with writes from one thread not immediately visible to all others; they do not provide sequentially consistent (SC) memory [8]. For some, such as Intel x86 [2] and IBM Power [1], the vendor documentation is in inevitably ambiguous informal prose, leading to confusion.

Recent progress has been made towards the formal specification of the behaviour of multiprocessors, e.g. [9, 5]. These formal foundations allow us to step towards verification of concurrent systems code running on top of deployed multiprocessors e.g. [3, 4].

But much remains to be done. A classical path to verify software is to use a program logic, such as Hoare logic [7]. Recent work by Richard Bornat shows that such a logic is possible for weak memory [6].

Yet, this logic takes an unusual form, with the proofs being underpinned by a graph structure which describes the execution paths of the program. This unusual form asks all the more for tool support: we crucially need to develop algorithms and tools to automate the discovery of proofs in this new logic. This is what I would like Tristan to work on during his internship.

References

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