A slice of CATALAN combinatorics from the perspective of COXETER sortable elements of the symmetric group

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Summary

$1.\ \mathrm{COXETER}$ sortable elements of the symmetric group

- Symmetric group
- COXETER sortable elements

2. Non crossing partitions

- Geometric non crossing partitions
- Bijection with the *c*-sortable elements

3. Binary trees

- Binary trees and SSA algorithm
- Link with *c*-sortable elements

$1.\ \mathrm{COXETER}$ sortable elements of the symmetric group

- Symmetric group
- COXETER sortable elements
- 2. Non crossing partitions

3. Binary trees

Notations

Let $n \in \mathbb{N}^*$. We note \mathfrak{S}_n the symmetric group of order n.

An element $\sigma \in \mathfrak{S}_n$ is represented by its one line notation : $\sigma(1)\sigma(2)\ldots\sigma(n)$. Example : $\mathfrak{S}_3 = \{1 \ 2 \ 3, \ 2 \ 1 \ 3, \ 2 \ 3 \ 1, \ 1 \ 3 \ 2, \ 3 \ 1 \ 2, \ 3 \ 2 \ 1\}.$

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COXETER group :
$$\mathfrak{S}_n \simeq \begin{pmatrix} s_1, \dots, s_{n-1} \\ s_i s_j = s_j s_i \\ s_i s_{i+1} s_i = s_{i+1} s_i s_{i+1} \end{pmatrix}$$
 $\forall 1 \le i \le n-1$
 $s_i s_{i+1} s_i = s_{i+1} s_i s_{i+1}$ $\forall 1 \le i \le n-2$

COXETER elements

A COXETER element of a COXETER group W is an element of W that is the product of all the generators exactly once.

Examples : $s_1 s_2 ... s_{n-1}$, $s_1 s_3 ... s_{n-1} ... s_4 s_2$.

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A COXETER element of \mathfrak{S}_n is permutation c of \mathfrak{S}_n that is a great cycle of the form $(1, a_1, \ldots, a_k, n, b_l, \ldots, b_1)$ where k + l = n - 2, $a_1 < \cdots < a_k$ and $b_1 < \cdots < b_l$.

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It is characterized by a partition of $\{2, \ldots, n-1\} = \underbrace{\{a_1, \ldots, a_k\}}_{L_c} \sqcup \underbrace{\{\underline{b_1, \ldots, b_l}\}}_{R_c}$.

Examples : let n = 6, (1, 3, 4, 6, 5, 2), (1, 6, 5, 4, 3, 2), (1, 2, 3, 4, 5, 6), (1, 5, 2, 6, 3, 4).

Binary trees

$\operatorname{COXETER}$ sortable elements

Let c be a COXETER element of \mathfrak{S}_n . A permutation $\sigma \in \mathfrak{S}_n$ is c-sortable [Rea05] if its one line notation avoids the following patterns :

• $ki \dots j$ for i < j < k and $j \in L_c$ • $j \dots ki$ for i < j < k and $j \in R_c$.

Binary trees

COXETER sortable elements

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• $ki \dots j$ for i < j < k and $j \in L_c$ • $j \dots ki$ for i < j < k and $j \in R_c$.

Examples : let n = 7 and c = (1, 3, 4, 6, 7, 5, 2).

- 7416325 is not c-sortable because it contains a pattern $ki \dots j$ with $j = 6 \in L_c$.
- 6521743 is not c-sortable because it contains a pattern $j \dots ki$ with $j = 5 \in R_c$.
- 3167425 is *c*-sortable because it avoids all the patterns.

CATALAN

Let $n \in \mathbb{N}^*$. For any COXETER element c of \mathfrak{S}_n , there are $C_n = \frac{1}{n+1} \binom{2n}{n}$ *c*-sortable elements.

Binary trees

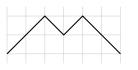
CATALAN

Let $n \in \mathbb{N}^*$. For any COXETER element c of \mathfrak{S}_n , there are $C_n = \frac{1}{n+1} \binom{2n}{n}$ c-sortable elements.

This means the c-sortable elements are in bijection with all the objects enumerated by the CATALAN numbers.

Examples :

 DYCK paths of length 2n



Triangulations of (n+2)-gons



Well parenthesized expressions of n+1 factors

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((a(bc))d)(ef)
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2. Non crossing partitions

- Geometric non crossing partitions
- Bijection with the *c*-sortable elements

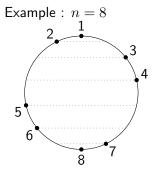
3. Binary trees

Labeling the circle

Let $n\in \mathbb{N}^*$ and place on a circle the numbers from 1 to n such that :

- 1 is at the highest point and n is at the lowest point,
- no two numbers are on the same height,
- reading from top to bottom the numbers are increasing.

Non crossing partitions



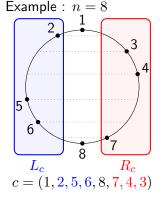


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A *c*-labeling is a labeling such that on the left are the elements of L_c and on the right the elements of R_c



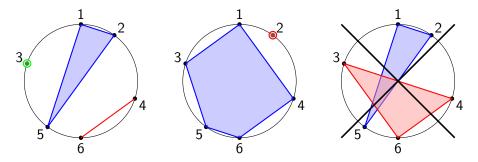
Non crossing partitions $\circ \circ \circ \circ \circ \circ$

Binary trees

Non crossing partitions

A *c*-non crossing partition is a set of non crossing polygons with vertices the marked points of a *c*-labeled circle. Single points and segments are considered as polygons.

Examples : n = 6 and c = (1, 3, 5, 6, 4, 2)



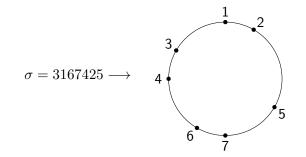
Non crossing partitions

Binary trees

Bijection with the *c*-sortable elements

Theorem (READING [Rea05])

The set of c-sortable elements is in bijection with the set of c-non crossing partitions via an explicit map called nc_c .



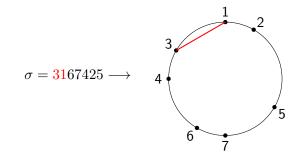
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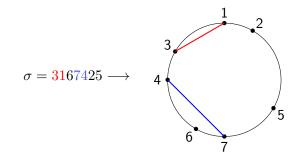
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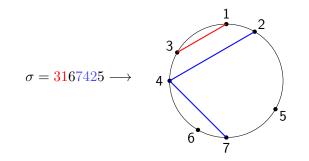
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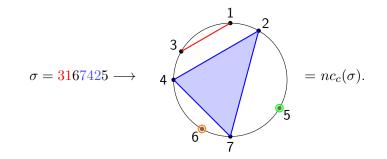
Non crossing partitions

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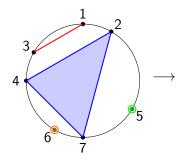
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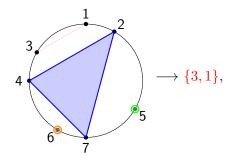
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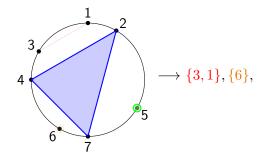
The inverse map of nc_c can be computed by selecting the polygons of a c-non crossing partition in a specific order [Gob18].



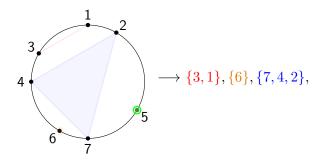
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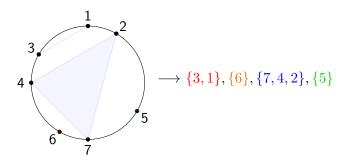
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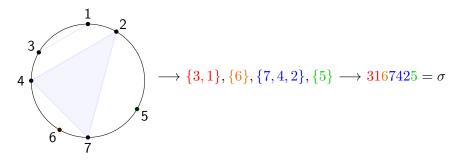
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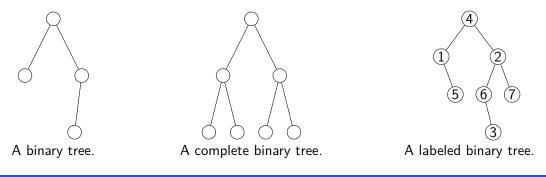
3. Binary trees

- Binary trees and SSA algorithm
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A binary tree is either an empty tree or a node with exactly one left child and one right child that are binary trees. The size of a binary tree is the number of nodes in the tree.

Examples :

Binary trees



 $\begin{array}{c} \text{COXETER sortable elements of the symmetric group} \\ \text{OOOOO} \end{array}$

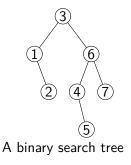
Non crossing partitions

Binary trees

Binary search trees, descending trees

A binary search tree is a labeled binary tree such that the label of each node is larger than the labels of its left child and smaller than the labels of its right child.

Examples :

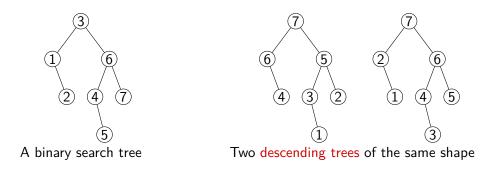


Binary trees

Binary search trees, descending trees

A binary search tree is a labeled binary tree such that the label of each node is larger than the labels of its left child and smaller than the labels of its right child. A descending tree is such that the label of each node is larger than the labels of its descendants.

Examples :



SSA algorithm

Non crossing partitions

Binary trees

Theorem (SSA algorithm [HNT04])

There is an explicit bijection

$$\mathfrak{S}_n \simeq \begin{cases} (T,Q) & T \text{ is a binary search tree of size } n \text{ and} \\ Q \text{ is a descending tree of the same shape as } T \end{cases}$$

Example : Let n = 7 and $\sigma = 2154763$.

$$T(\sigma) = Q(\sigma) =$$

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Example : Let n = 7 and $\sigma = 2154763$. (position = 7)

 $T(\sigma) = 3$ $Q(\sigma) = 7$

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Example : Let n = 7 and $\sigma = 2154763$. (position = 6)

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SSA algorithm

Non crossing partitions

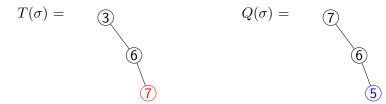
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Example : Let n = 7 and $\sigma = 2154763$. (position = 5)



SSA algorithm

Non crossing partitions

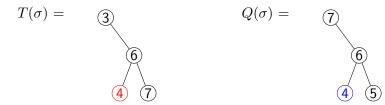
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Example : Let n = 7 and $\sigma = 2154763$. (position = 4)



SSA algorithm

Non crossing partitions

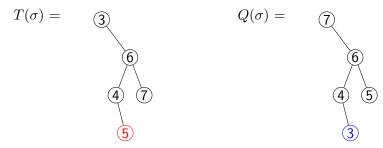
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SSA algorithm

Non crossing partitions

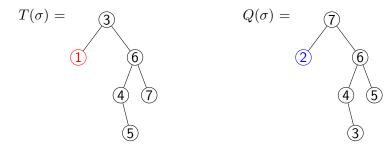
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Example : Let n = 7 and $\sigma = 2154763$. (position = 2)



SSA algorithm

Non crossing partitions

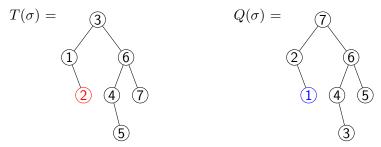
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Example : Let n = 7 and $\sigma = 2154763$. (position = 1)



 $\begin{array}{c} \text{COXETER sortable elements of the symmetric group} \\ \text{OOOOO} \end{array}$

SSA algorithm

Non crossing partitions

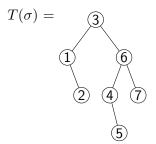
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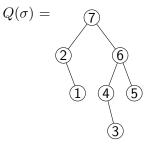
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Example : Let n = 7 and $\sigma = 2154763$.





Non crossing partitions

Binary trees 0000●0

Sylvester congruence

SSA algorithm : $\sigma \in \mathfrak{S}_n$ is encoded by $(T(\sigma), Q(\sigma))$.

What happens if we forget $Q(\sigma)$? Can we describe all $\sigma' \in \mathfrak{S}_n$ s.t. $T(\sigma) = T(\sigma')$?

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SSA algorithm : $\sigma \in \mathfrak{S}_n$ is encoded by $(T(\sigma), Q(\sigma))$.

What happens if we forget $Q(\sigma)$? Can we describe all $\sigma' \in \mathfrak{S}_n$ s.t. $T(\sigma) = T(\sigma')$?

Yes! $T(\sigma) = T(\sigma')$ iff σ' can be obtained from σ by a series of transformations of the form $ki \dots j \leftrightarrow ik \dots j$ with $i < j < k \longrightarrow$ Sylvester congruence on \mathfrak{S}_n .

Example : $\sigma' = 5421763$ has the same binary search tree than $\sigma = 2154763$

 $2154763 \rightarrow 2514763 \rightarrow 2541763 \rightarrow 5241763 \rightarrow 5421763$

Non crossing partitions

Binary trees ○○○○○●

Link with *c*-sortable elements

If $c = (1, 2, 3, \dots, n-1, n)$, then we have the following one to one maps :

2143

 σ with no $ki \dots j$ pattern = a *c*-sortable element

Non crossing partitions

Binary trees ○○○○○●

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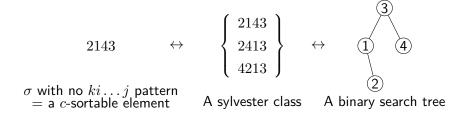
$$2143 \qquad \leftrightarrow \qquad \left\{ \begin{array}{c} 2143\\ 2413\\ 4213 \end{array} \right\}$$

Non crossing partitions

Binary trees

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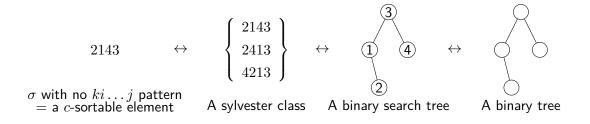


Non crossing partitions

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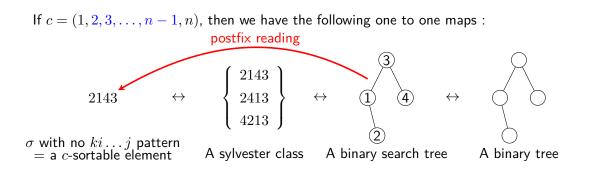
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Non crossing partitions

Binary trees ○○○○○●

Link with *c*-sortable elements



The map from binary trees to *c*-sortable elements can be directly obtained with a postfix reading of the associated binary search tree.

READING's map nc_c : {*c*-sortable elements} \rightarrow {*c*-non crossing partitions} is well defined for any (finite rank) COXETER group, and is a bijection in all finite COXETER groups.

In infinite COXETER groups, it is only injective, but never surjective. For example, in type \tilde{A}_1 and c = st, all reflections are *c*-non crossing but all the following ones are not in the image of nc_c : tst, tstst, tststst, ...

My goal is to define a generalized notion of c-sortable elements, at least in the affine types, such that READING's map can be naturally extended to be a bijection.

Thank you for your attention!

• What about other COXETER elements? \longrightarrow another tree structure : Cambrian trees [CP17].

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• What about other finite COXETER groups? \longrightarrow *W*-CATALAN numbers [Rei97].

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• What about other finite COXETER groups? \longrightarrow W-CATALAN numbers [Rei97].

What about infinite COXETER groups? → It's complicated... Non crossing partitions are well defined (in a more algebraic way) as well as READING bijection, but it never is surjective. In the affine type, there's a hope to generalize the *c*-sortable elements to a larger family that becomes in bijection with the non-crossing partitions → my thesis.

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