## Evaluating regulation policies for subways with model checking

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## Introduction

Every day, millions of people take the subway !

## Challenges:

- Prove the safety of subway networks
- Ensure the efficiency of subways regarding delays with a regulation policy



## Introduction: related work

## Safety

Model checking is used to prove the security of critical sections ${ }^{1}$ (e.g. signaling system)

## Efficiency

Simulation of the physic reality of subways from a very specific situation ${ }^{2}$

[^0]
## Introduction: our approach

## Hypothesis

The safety of the subway networks we study is ensured.

- Model checking offers formal guarantees
- It can be used to evaluate efficiency of subways
- That is : evaluating regulation policies
- Use the model checker PRISM ${ }^{3}$

[^1]
## Outline

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## Required features in the model

We need a formal model to represent subway networks.

## Need for randomness

Delays are unpredictable and conveniently represented through probabilities.

## Need for nondeterminism

Regulation policies can increase or decrease the dwell time of subways in station. This can be seen as nondeterminism.

## Model: Markov Decision Process (MDP)



An example of an MDP

## Regulation policy

- Chooses the behavior of trains according to the state of the system
- Resolves the nondeterminism of the MDP


## Our Goal

Design regulation policies and evaluate their efficiency at recovering from a delay

## Glasgow: a simple topology

- Real systems are often too complex for formalization
- Simpler the system, simpler the model!
- First, we study a ring system



## Outline

## Space and Time Discretization



## Parameters of interest

- Time discretization step: $\Delta t$
- Space discretization step: $\Delta d$
- Probabilities: $p_{a}, p_{b}$ (also $q_{a}=1-p_{a}$ and $q_{b}=1-p_{b}$ )
- Number of intermediate steps: $k$
- Number of trains: $n b_{\text {train }}$



## Choosing the probabilities: data from Santiago

Distribution of time to go from one station to another (7090 samples)


## Choosing the probabilities: data from Santiago

Distribution of time to go from one station to another (7090 samples)


## Choosing the parameters: from the Glasgow subway

## Choosing $n b_{\text {train }}$ :

- Peak time: $n b_{\text {train }}=6$
- Off-peak time: $n b_{\text {train }}=4$

We have the following relation between $k, \Delta t$ and $p_{a}$ :

$$
k \times \Delta t=p_{a} \times 66 s
$$

- $k=5$
- $k=10$
- $\Delta t \simeq 10 s$
- $\Delta d=140 m$
- $\Delta t \simeq 5 s$
- $\Delta d=70 m$


## How to estimate delay?

- empty station
: subway in station
: subway of interest


9
8 7

$$
\alpha=\frac{d(\text { current }, \text { next })}{d(\text { previous }, \text { current })+d(\text { current }, \text { next })} \in[0,1]
$$

## How to estimate delay?

- empty station


Delay: $\alpha \notin[0.4,0.6]$

## Extreme cases

- empty station
: subway in station
: subway of interest



## A simple regulation policy

Chooses the dwell time in station as a function of $\alpha$ :


## Properties to be checked

## Safety property

Two trains must not collide: $P_{\max =0}(G \neg$ "collision" $)$

Efficiency of the regulation policy given an initial configuration

- Recovering time from an unbalanced configuration: $P_{\text {min }=?}\left(F_{\leq n} \neg "\right.$ delay" $)$
- Avoiding delays from a balanced configuration: $P_{\text {max }=?}\left(F_{\leq n}\right.$ "delay" $)$


## Outline

## First attempt

- Automated generation of prism models and properties on which prism may work

- Prism : unable to build the state space for $n b_{\text {train }}=4, k=5$ (smaller model of interest), the properties cannot be verified


## Abstraction: reduce the size of the model

: empty station
: subway in station


## Abstraction: station ids are irrelevant

- empty station
: subway in station



## Abstraction of our model: description



- point: distance between a train and its previous station
- nb_station: number of stations between a train and its successor


## Abstraction of our model: some results

| Model <br> Number of trains | Before <br> abstraction | After <br> abstraction |
| :--- | :---: | :---: |
| 3 trains | $2.1 \times 10^{8}$ states | $3.5 \times 10^{5}$ states |
|  | $3.4 \times 10^{9}$ transitions | $8.3 \times 10^{5}$ transitions |
| 4 trains | Not built | $2.0 \times 10^{7}$ states |
|  | in PRISM | $5.7 \times 10^{7}$ transitions |

Table: Size of the model in terms of number of states and transitions

## Soundness of the model

## The model must satisfy the safety property!

- Provable in Prism with four trains
- Prism cannot build the model with six trains

A new abstraction:

- A simpler model: encompasses the previous one
- Every transition becomes nondeterministic
- Safety property was proven with 6 trains


## Evaluating efficiency: $n b_{\text {train }}=4, k=5$



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Probabilty to recover from a delay with nb_train $=4$ and $k=5$ as a function of time


## Evaluating efficiency: $n b_{\text {train }}=4, k=10$



## Evaluating efficiency: $n b_{\text {train }}=4, k=10$



## Evaluating efficiency: $n b_{\text {train }}=4, k=10$

Probabilty to recover from a delay with nb_train $=4$ and $k=10$ as a function of time


## Evaluating efficiency: $n b_{\text {train }}=6, k=10$

Probabilty to recover from a delay with nb_train $=6$ and $k=10$ as a function of time


## Future Work

- Assess more accurately the efficiency of the regulation policy
- Refine the abstraction of the model
- Study another modelisation of the speed of subways
- What about a new definition of delay ?


## Discrete Time Markov Chain (DTMC)



## PCTL logic

The PCTL ${ }^{45}$ logic uses sevral connectors:

- The usual connectors of propositional logic
- Temporal connectors :
- Next : X $\phi$
- Eventually: $F \phi$
- Bounded eventually : $F^{\leq n} \phi$
- A probabilistic connector : $P_{\alpha p}$ with $\alpha \in\{\leq,<, \geq,>\}$ and $p \in[0,1]$

[^2]
## PCTL logic

- Two trains never collide :

$$
\phi=P_{\leq 0}\left(F \phi_{\text {collision }}\right)
$$

- If a train has some delays it will catch it up within 10 steps with a high probability:

$$
\phi=\phi_{\text {delay }} \Rightarrow P_{\geq 0.9}\left(F \leq 10 \neg \phi_{\text {delay }}\right)
$$

## Choosing the probabilities: data from Santiago

Distribution of time to go from one station to another (7090 samples)


## Parameters 2

Data collected from actual subway rail system:

- Total duration of the course in Glasgow : $t_{\text {tot }}=24 \mathrm{~min}$
- Length of a complete circuit in Glasgow : $d=10.5 \mathrm{~km}$
- Usual speed of subways : $v$ between 30 and $40 \mathrm{~km} . \mathrm{h}^{-1}$
- Restriction on the probability : $p \geq 0.8$


[^0]:    ${ }^{1}$ Automated verification and validation of signaling systems in PTC and CBTC environements, Smith et al., 2012
    ${ }^{2}$ Railroad simulation using opentrack. A. Nash and D. Huerlimann, 2004

[^1]:    ${ }^{3}$ PRISM 4.0: Verification of Probabilistic Real-time Systems, Kwiatkowska et al., 2011

[^2]:    ${ }^{4}$ A logic for reasoning about time and reliability, Hanson et al., 1994
    ${ }^{5}$ Automatic Verification of Finite-state Concurrent Systems Using Temporal Logic Specifications, Clarke et al., 1986

