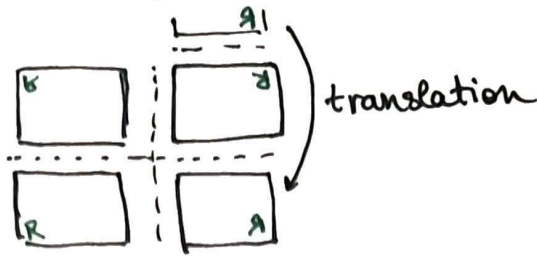


Translation surfaces: definitions

Recall last time:



1. Constructive definition

P : a finite set of polygons in \mathbb{R}^2 (maybe non convex)

$E(P)$: the edges of the polygons in P

For every $e \in E(P)$, we define \underline{n}_e the unit vector which is orthogonal to e and points toward the interior of the polygon $e \in P$



A map $f: E(P) \rightarrow E(P)$ is a pairing if

1. $f \circ f = \text{id}$, f has no fixed point
2. If $e_2 = f(e_1)$, then there exist $w_1 \in \mathbb{R}^2$ s.t. $e_2 = e_1 + w_1$
we denote f_{e_1} the translation by w_1 in \mathbb{R}^2
3. then $f_{e_1}(\underline{n}_{e_1}) = -\underline{n}_{e_2}$

Ex: $e_3 \xrightarrow{f} e_1$ We could not have $f(e_1)$ because $f_{e_1}(\underline{n}_{e_1}) = \underline{n}_{e_3}$

We define $X = \bigsqcup_{P_i \in P} P_i / \sim_f$

where

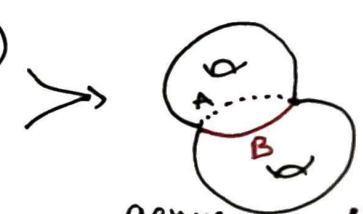
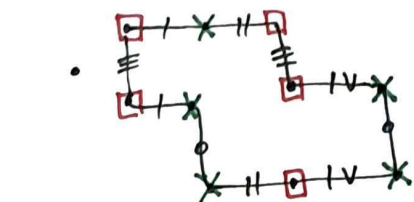
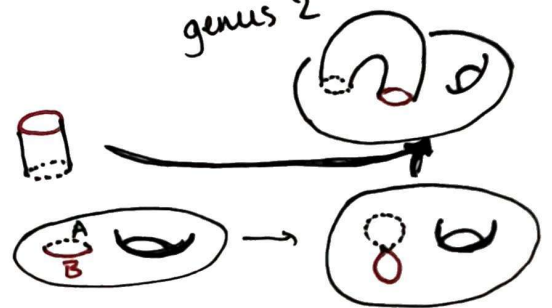
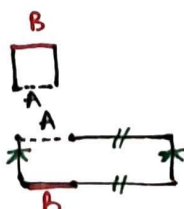
$\sim_f: x_1 \sim x_2 \iff x_1 \in e_1, x_2 \in e_2, f(e_1) = f(e_2) \text{ and } f_{e_1}(x_1) = x_2$

Call X a translation surface obtained by P and f if X is connected.

Ex: $\left\{ \begin{matrix} \text{[diagram 1]} \\ \text{[diagram 2]} \end{matrix} \right\} = P$, f given by the marking $0, 1, 5, 1, \dots$

The torus is a translation surface obtained by P and f
 $X = \bigsqcup_{P_i \in P} P_i / \sim_f$

$P = \{P_0\}$

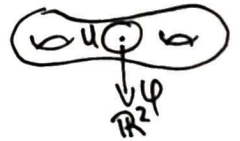


The vertices of the polygon don't identify to one point, but two.

Rmk: If $f: A \rightarrow A$ with $A \subseteq E(P)$ satisfies 1. 2. and 2' then we can define a translation surface with boundary.

Claim: Given any rational polygon P (the angle of P are rational multiple of π), there exist finitely many reflections of P along its sides (produced by the unfolding process), and there exists a pairing f s.t. $X = \bigcup_{P_i \text{ reflection of } P} P_i / \sim_f$ is a translation surf.
(maybe proved later)

Why is X (obtained by P and f) a surface?



Local geometry of X

The map $\bigcup_{P_i \in P} P_i \xrightarrow{\pi} X = \bigcup P_i / \sim_f$ is such that:

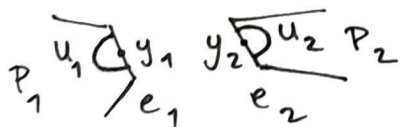
1. $\pi|_{\bigcup P_i}$ is injective
2. $\pi|_{\bigcup_{e \in E(P)} e}$ is 2 to 1
3. $\pi|_{\bigcup_{v \in V(P)} v}$ is many to 1 (where $V(P)$ is the set of vertices)

In case 1., for $x \in X$, $y \in P_i$ s.t. $\pi(y) = x$, we have $U_y \subset P_i$ a Euclidean ball $U_x = \pi(U_y)$



$(U_x, \pi^{-1}|_{U_x})$ is a chart around x .

In case 2., $x = \pi(y_1) = \pi(y_2)$, $y_i \in e_i$



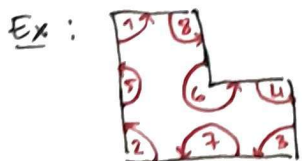
We have half-disks U_1 and U_2 centered at y_1 and y_2 within P_1 and P_2 .



$$U_x = \pi(U_1) \cup \pi(U_2)$$

$U_x \xrightarrow{\varphi} \mathbb{D}_{y_1}$ gives a chart around x .

In case 3., $x = \pi(y_1) = \dots = \pi(y_k)$

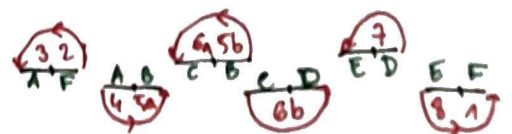


Topologically, they are glued to a disk



This shows that X is a surface.

We can also use the half-disks construction:



Claim Let X be obtained by P and f
 $x \in X, x = \pi(y_1) = \dots = \pi(y_k)$ for $y_i \in V(P)$
 Then there is a neighborhood U_x of x obtained by the half-disks construction for $2m$ half-disks, $m \in \mathbb{N}_{>0}$.

Proof: Every identification is by translation
 so we must do complete "turns" before "closing" the neighborhood.
 The total angle has to be $m \times 2\pi$. \square



Remark: The cone angle at x is the sum of cone angles of y_i 's in their polygon P_i . The cone angle is $m \cdot 2\pi$ with $m \in \mathbb{N}_{>0}$.

Computation of the genus of X

$$2 - 2g = \chi(X) = \#V - \#E + \#F = \#V - \frac{\#E(P)}{2} + \#P$$

$\triangle \#V \neq \#V(P)$

Let x_1, \dots, x_s be points in X from vertices,
 and $m_1 \cdot 2\pi, \dots, m_s \cdot 2\pi$ be their cone angles.

Define $n_i = m_i - 1$ the degree.

Lemma: $\boxed{\sum_{i=1}^s n_i = 2g - 2}$

PP: ~~Use~~ ~~when~~ P is a polygon with l edges,
 The total angle (sum of all inner angles) of P is $(l-2)\pi$

We compute:

$$2\pi \sum_{i=1}^s n_i + 2\pi \#V = \sum_{i=1}^s (n_i + 1) 2\pi = \sum_{i=1}^s m_i 2\pi = \sum_{P_j \in P} \text{total angle of } P_j$$

$$= \sum_{P_j \in P} (l_{P_j} - 2)\pi = \pi \#E(P) - 2\pi \#P$$

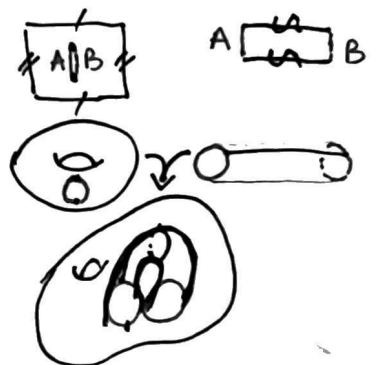
So: $\sum_{i=1}^s n_i = -\#V + \frac{\#E(P)}{2} - \#P = -\chi(S) = 2g - 2$ \square

Remark 1: $\forall i, m_i \geq 1$ so $n_i \geq 0$ and there is no translation surface of genus 0.

Remark 2: $\forall g \geq 1$, there is a translation surface of genus g .

We know ~~the~~ of genus 1.

We can add genus with a slit construction:



With k slits (and k rectangles) on a square,
 we get a ~~the~~ surface of genus $k+1$.